Adaptive Sensing for Recovering Structured Sparse Sets

Ervin Tánczos, Rui Castro Eindhoven University of Technology

Structures Seminar 16.10.2015



Motivation

We are interested in recovering the support (or detecting the presence) of an unknown signal.



Classical Framework

Let $\mathbf{x} = (x_1, \dots, x_n)^T \in \mathbb{R}^n$ denote the unknown signal where $x_i = \begin{cases} \mu & , \text{ if } i \in S \\ 0 & , \text{ if } i \notin S \end{cases}$,

with $\mu > 0$ fixed and $S \in \mathcal{C}$ where \mathcal{C} is a class of sets.



Classical Framework

We observe

$$Y_i = x_i + W_i, \ W_i \stackrel{iid}{\sim} N(0,1), \ i = 1, \dots, n .$$

Our goal is to recover S or to detect its presence (decide between H_0 : $S = \emptyset$ and H_1 : $\emptyset \neq S \in C$).

How does μ need to scale so that the above tasks are possible?

Classical Framework

Depends on the class C. From now on, assume $|S| = s \ll n \ \forall S \in C$ (sparse signals). We want $\max_{\substack{S \in C}} \mathbb{P}(Error)$ to be small.

Non-Adaptive	Detection	Recovery
s-sets	$\sqrt{\log n}$	$\sqrt{\log n}$
unions of k disjoint s -intervals	$\sqrt{\frac{1}{s}\log n}$	$\sqrt{\frac{1}{s}\log n}$
unions of k disjoint s -stars	$\sqrt{\log n}$	$\sqrt{\log n}$
$\sqrt{s} \times \sqrt{s}$ submatrices	$\sqrt{\frac{1}{\sqrt{s}}\log n}$	$\sqrt{\frac{1}{\sqrt{s}}\log n}$

 $\max_{S \in \mathcal{C}} \sum_{i \in S} X_i \text{ does the job (in the sparse regime).}$

See e.g. Lugosi et. al. (2010): On combinatorial testing problems; Arias-Castro et. al. (2011): Detection of an Anomalous Cluster in a Network

・ロッ ・雪 ・ ・ ヨ ・ ・ ヨ ・

э

Learning to learn



- How can we take advantage of feedback?
- How much can we gain?

ション ふゆ アメリア メリア しょうくの

Framework

The unknown signal and the goals are the same as before. Measurement model:

$$Y_t = x_{A_t} + W_t, \ W_t \stackrel{iid}{\sim} N(0,1), \ t = 1,\ldots,n$$

where $A_t \in \{1, ..., n\}$ can depend on past observations $\{A_j, Y_j\}_{j=1}^{t-1}$.

ション ふゆ アメリア メリア しょうくの

Framework

The unknown signal and the goals are the same as before. Measurement model:

$$Y_t = x_{A_t} + \Gamma_t^{-1/2} W_t, \ W_t \stackrel{iid}{\sim} N(0,1), \ t = 1, 2, \dots,$$

where $A_t \in \{1, ..., n\}, \Gamma_t > 0$ can depend on past observations $\{A_j, \Gamma_j, Y_j\}_{j=1}^{t-1}$, and must satisfy

$$\mathbb{E}_{S}\left(\sum_{t} \Gamma_{t}\right) \leq n, \ \forall S \in \mathcal{C} \ .$$

・ロト ・雪 ・ ミート ・ ヨー うらつ

Simple procedure for recovery

Let C be the class of all *s*-sparse sets and suppose we wish to recover the support (we want \widehat{S} s.t. $\max_{S \in C} \mathbb{P}_{S}(\widehat{S} \neq S) \leq \varepsilon$).

Algorithm

- Fix $\Gamma_t = \Gamma = 1/3 \ \forall t \in \mathbb{N}$
- For each entry x_i , i = 1, ..., n do the following:
 - Measure $Y_{i,j} = x_i + \Gamma^{-1/2} W_i$, $j = 1, \dots, \tau_i$, where $\tau_i = \min\{j : Y_{i,j} \le 0\} \land \log_2(n/\varepsilon)$.
- $i \in \widehat{S} \iff Y_{i,j} > 0 \forall j = 1, \dots, \log_2(n/\varepsilon).$

▲□▶ ▲圖▶ ▲臣▶ ★臣▶ ―臣 …の�?

Simple procedure for recovery - analysis

For $i \notin S$

$$\mathbb{P}(i \in \widehat{S}) = \mathbb{P}(Y_{i,j} > 0 \,\, orall j) \leq (1/2)^{\log_2(n/arepsilon)} = arepsilon / n \,\,.$$

For $i \in S$

$$\mathbb{P}(i \notin \widehat{S}) \leq \mathbb{P}(\exists j : Y_{i,j} \leq 0) \leq \frac{\log_2(n/\varepsilon)}{2} e^{-\mu^2/6} \leq \varepsilon/s$$

whenever
$$\mu \geq \sqrt{6\left(\log \frac{s}{\varepsilon} + \log \frac{\log_2(n/\varepsilon)}{2}\right)}.$$

Hence

$$\mathbb{P}_{\mathcal{S}}(\widehat{S} \neq \mathcal{S}) \leq \sum_{i \notin \mathcal{S}} \mathbb{P}(i \in \widehat{S}) + \sum_{i \in \mathcal{S}} \mathbb{P}(i \notin \widehat{S}) \leq \varepsilon \;.$$

・ロト ・ 日 ・ ・ 日 ・ ・ 日 ・ ・ つ へ ()

Simple procedure for recovery - analysis

How much precision do we use in expectation?

$$\mathbb{E}_{S}\left(\sum_{t} \Gamma_{t}\right) \leq \Gamma\left(\sum_{i \notin S} 2 + \sum_{i \in S} \log_{2}(n/\varepsilon)\right) \leq \frac{1}{3}(2n + s\log(n/\varepsilon)) \leq n$$

if $s \ll n$.

To summarize, this simple procedure succeeds when

$$\mu\gtrsim \sqrt{\log s + \log\log n + \log rac{1}{arepsilon}}$$
 .

Reminder - SLRT (Wald)

We wish to test H_0 : $Y_j \sim N(0, \Gamma^{-1})$ and H_1 : $Y_j \sim N(\mu, \Gamma^{-1})$, $j \in \mathbb{N}$ with as few observations as possible (in expectation) with prescribed error probabilities α, β . Consider the process

$$Z_0 = 0, \ Z_t = \sum_{j=1}^t \log rac{f_1(Y_i)}{f_0(Y_i)}, t = 1, 2, \dots$$

Let $I = \log \beta < 0 < u = \log(1/\alpha)$ and $T = \inf\{t : Z_t \notin (I, u)\}$. We then have $\mathbb{P}_0(Z_T \ge u) \le \alpha$ and $\mathbb{P}_1(Z_T \le I) \le \beta$.

As $\Gamma \to 0$ we also have

•
$$\mathbb{P}_0(Z_T \ge u) \to \alpha \text{ and } \mathbb{P}_1(Z_T \le l) \to \beta$$

•
$$E_0(T) \approx rac{2}{\Gamma\mu^2} \log rac{1}{eta}$$
 and $E_1(T) \approx rac{2}{\Gamma\mu^2} \log rac{1}{lpha}$

Refinement

Replace the core of the previous procedure with a SLRT to test between $x_i = 0$ and $x_i = \mu$. Set Type I and II error probabilities to be $\alpha = \varepsilon/n$ and $\beta = \varepsilon/s$. We have $\mathbb{P}_S(\widehat{S} \neq S) \le \varepsilon$ as before.

The precision used (in expectation) is

$$\mathbb{E}_{\mathcal{S}}\left(\sum_{t} \Gamma_{t}\right) \leq \frac{2}{\mu^{2}}\left(n\log\frac{s}{\varepsilon} + s\log\frac{n}{\varepsilon}\right)$$

If *n* is large (and $s \ll n$) this is at most *n* whenever

$$\mu \geq \sqrt{2\log rac{s}{arepsilon} + o(1)} \;.$$

This is optimal.

ション ふゆ アメリア メリア しょうくの

Detection

What about detection? Easy: set α as before and $\beta = \sqrt[s]{\varepsilon}$. This ensures that at least one signal component is found w.p. $1 - \varepsilon$ under the alternative.

Adaptive	Detection	Recovery
s-sets	$\sqrt{1/s}$	$\sqrt{\log s}$
unions of k disjoint s -intervals	$\sqrt{1/s}$?
unions of k disjoint s -stars	$\sqrt{1/s}$?
$\sqrt{s} \times \sqrt{s}$ submatrices	$\sqrt{1/s}$?

Scaling laws do not depend on the structure anymore (as long as we have symmetry in the class)!

Structured Recovery



Structured Recovery



Structured Recovery



Structured Recovery



Structured Recovery



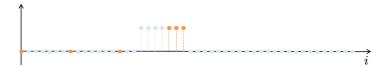
▲□▶ ▲□▶ ▲□▶ ▲□▶ ▲□ ● ● ●

Structured Recovery

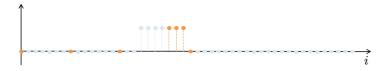


▲□▶ ▲□▶ ▲□▶ ▲□▶ ▲□ ● ● ●

Structured Recovery



Structured Recovery



Structured Recovery

Main idea: take a "noiseless case" algorithm for support recovery and "robustify" it against noise by using SLRTs.

Typically the algorithm will have two phases:

- Search: Find an active component (can also use random search)
- Refinement: Exploit structure around that component

Algorithms may alternate between the two phases (for instance in case of unions of stars).

The main difference between the two phases is that the error probabilities for the SLRTs are set differently.

・ロト ・ 日 ・ ・ 日 ・ ・ 日 ・ ・ つ へ ()

Detection

Still considering probability of error we get (recall we are in the sparse regime $s \ll n$).

Adaptive	Detection	Recovery
s-sets	$\sqrt{1/s}$	$\sqrt{\log s}$
unions of k disjoint s -intervals	$\sqrt{1/s}$	$\sqrt{\frac{\log k}{s}}$
unions of k disjoint s -stars	$\sqrt{1/s}$	$\sqrt{\frac{\log k}{s}}$
$\sqrt{s} \times \sqrt{s}$ submatrices	$\sqrt{1/s}$	$\sqrt{1/s}$

Detection

For technical reasons we only managed to show lower bounds for the recovery problem considering $\max_{\substack{S \in \mathcal{C}}} \mathbb{E}_{S}(|\widehat{S} \triangle S|) \leq \varepsilon$.

Adaptive	Detection	Recovery
s-sets	$\sqrt{1/s}$	$\sqrt{\log s}$
unions of k disjoint s -intervals	$\sqrt{1/s}$	$\sqrt{\frac{\log(ks)}{s}}$
unions of k disjoint s -stars	$\sqrt{1/s}$	$\sqrt{\frac{\log(ks)}{s}}$
$\sqrt{s} \times \sqrt{s}$ submatrices	$\sqrt{1/s}$	$\sqrt{\frac{\log s}{s}}$

Adaptive algorithms can improve on non-adaptive ones by

- Better mitigating the effects of noise $\log n \rightarrow \log s$
- Better capitalizing on structure (in certain cases) $\sim 1/s$

▲□▶ ▲□▶ ▲□▶ ▲□▶ ▲□ ● ● ●

Framework

The unknown signal and the goals are the same as before.

Different measurement model:

$$Y_t = \langle x, A^{(t)} \rangle + W_t, W_t \stackrel{iid}{\sim} N(0, 1), t = 1, 2, \dots,$$

where $A^{(t)} \in \mathbb{R}^n$ can depend on past observations $\{A^{(j)}, Y_j\}_{j=1}^{t-1}$,
and must satisfy

$$\mathbb{E}_{S}\left(\sum_{t} \|A^{(t)}\|_{F}^{2}\right) \leq n, \ \forall S \in \mathcal{C} \ .$$

ション ふゆ アメリア メリア しょうくの

Detection

Consider the energy test $Y_1 = \langle x, 1 \rangle + W_1$, where $1 \in \mathbb{R}^n$ is a vector of ones and $\Psi = 1\{Y_1 > s\mu/2\}$.

We have

$$\max_{i=0,1} \mathbb{P}_i(\Psi \neq i) \leq \varepsilon ,$$

whenever $\mu \ge \sqrt{\frac{8}{s^2} \log \frac{1}{2\varepsilon}}$. This is optimal among all tests (adaptive or non-adaptive).

Structure and adaptivity do not play a role.

Arias-Castro (2012): Detecting a vector based on linear measurements

Simple procedure for recovery

Consider the 1-sparse case, and a binary search algorithm.

Let $A^{(1)} \in \mathbb{R}^n$ s.t. $A^{(1)}_i = \mathbf{1}\{i \le n/2\}$ and $Y_1 = \langle x, A^{(1)} \rangle + W_1$. If $Y_1 > \mu/2$ "go left" otherwise "go right", and iterate. This simple procedure has $\max_{S \in \mathcal{C}} \mathbb{P}_S(\widehat{S} \neq S) \le \varepsilon$ whenever

$$\mu \geq \sqrt{8\left(\log rac{\log_2 n}{2} + \log rac{1}{arepsilon}
ight)} \;.$$

(also $\sum \|A^{(t)}\|_F^2 \leq n$)

Similarly as before, replacing the observations by SLRTs (multiple measurements with small sensing energy) we can get rid of the log log term.

Recovery

One can use the insights gained above for structured recovery. For *s*-sparse sets do *s* binary searches in parallel.

For structured sets do two phases as before. In the search phase

- Intervals: Search for a block of activation.
- Stars: Search for the center of the active star.
- **Submatrices:** Search for rows that contain activation.

The refinement phases are "easy" when $s \ll n$ (compared to the search phases).

Malloy, Nowak (2013): Near-Optimal adaptive compressed sensing

Recovery

Recovery	Non-adaptive	Adaptive
s-sets	$\sqrt{\log n}$	$\sqrt{\log s}$
unions of k disjoint s -intervals	$\sqrt{\frac{\log n}{s^2}}$	$\sqrt{\frac{\log(ks)}{s^2}}$
unions of k disjoint s -stars	$\sqrt{\log n}$	$\sqrt{\frac{\log(ks)}{s^2}}$
$\sqrt{s} \times \sqrt{s}$ submatrices	$\sqrt{\frac{\log n}{\sqrt{s}}}$	$\sqrt{\frac{\log(ks)}{s}}$

Non-adaptive rates are necessary, adaptive ones are sufficient and except for submatrices also necessary.

Similar behavior as before.

Remark - number of measurements

Appeal of compressive sensing: few measurements ($\approx s \log n$). We lose this in the algorithms above.

Note that in binary search $||A^{(t)}||_0 = 2^{-t}$. This allows us to choose $||A^{(t)}||_F^2 \sim t2^{-t}$ and still satisfies $\sum_t ||A^{(t)}||_F^2 \leq n$. This way we get rid of the log log term (at the price of an increase in the constant).

Same can be done to all other algorithms \rightsquigarrow same performance, small number of measurements.

In the non-adaptive case *s* log *n* measurements are optimal. In the non-adaptive case we don't know (yet).

ション ふゆ アメリア メリア しょうくしゃ

Final remark

The crux of all adaptive sensing algorithms is the sampling strategy.

We aim to collect the most "informative" samples based on what we already learned.

Would a sampling strategy that at time t = 1, 2, ... decides what to do based on the posterior of $S|Y_1, ..., Y_{t-1}$ make sense?